

A neighborhood proximity based algorithm
for overlapping community structure
detection in weighted networks

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Problems & Ideas

The **neighborhood proximity** of a pair (u, v) is

$$\rho(u, v) = \left(\frac{\sum_{w, x \in N_{uv}} a_{uw} + a_{vw} + a_{wx}}{\sum_{w \in N_{uv}} s_w} \right) \left(\frac{\sum_{w \in N_{uv}} \min(a_{uw}, a_{vw}) + a_{uv}}{\min(s_u, s_v)} \right)$$

if $N_{uv} \neq \emptyset$. Otherwise, $\rho(u, v) = a_{uv} / \min(s_u, s_v)$.

The **weighted overlapping modularity** of a cover \mathcal{K} is

$$Q_{\text{wo}}(\mathcal{K}) = \frac{1}{|\mathcal{K}|} \sum_{C \in \mathcal{K}} \left(1 - \frac{\sum_{v \in C} |m_v|}{|C| \cdot |\mathcal{K}|} \right) \frac{d^{\text{int}}(C)}{d(C)},$$

and $m_v =$ no. of communities v belongs to.

- \mathcal{K} is optimal \iff
- 1) $d^{\text{int}}(C)/d(C) = 1, \forall C \in \mathcal{K}$
 - 2) $|\mathcal{K}|$ is highest possible.
 - 3) $|m_v|$ is least possible, $\forall v \in V$

Main Contributions

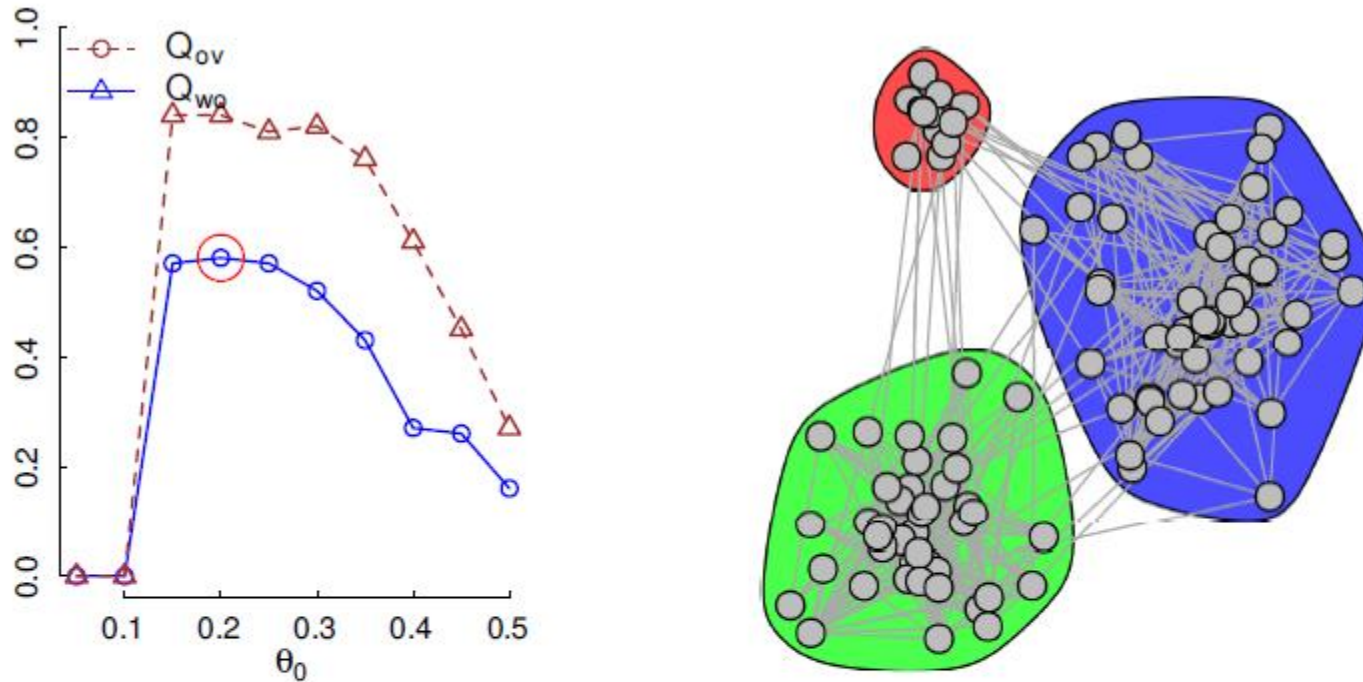


Figure: Community structure in US-POLBOOKS: The left panel plots the maximum values of Q_{ov} and Q_{wg} over 10 iterations of NPCD for each value of θ_0 in the range $\{0.05, 0.10, \dots, 0.50\}$. The right panel plots the community structure corresponding to the maximum value of Q_{wg} over the entire range of θ_0 .