

Group Relational Privacy Protection on Time-constrained Point of Interests

**Bo NING, Xiaonan Li, Fan YANG, Yunhao SUN,
Guanyu LI, George Y. YUAN**

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Model and Our Approach

- **$k^{m,n}$ -Anonymity Model:** Given a time-constrained data set $D(U, I)$ of users, a $k^{m,n}$ -anonymity is used to make the number of each vulnerable PoIs in D not less than k , where m and n denote the number of common PoIs of two-user and multi-user as the prior knowledge known by attackers, respectively.
- **Our Approach:** Anonymous Algorithm of User Relationship

Algorithm 1: Anonymous Algorithm of User Relationship

Input: time-constrained data set of user PoIs $D(U, I)$,
prior knowledge m, n, k , distance matrix M
Output: anonymous data set \mathcal{D}

```
1 for  $u, u' \in U$  do
2    $R_{single}^m \leftarrow R_{single}^m(u, u')$ ;
3 for  $u \in U$  and  $r \in R_{single}^m$  do
4    $R_{group}^n \leftarrow R_{group}^n(u, u')$ ;
5    $Tree \leftarrow Root(r) \cup Node(u, r)$ ;
6 if  $m \leq n$  then
7   anonymize  $R_{single}^m$  on  $D$  limited by  $M$ ;
8 else
9   for  $r \in R_{single}^m$  do
10    anonymize  $r$  on  $D$  limited by  $M$ ;
11    iteratively update and anonymize nodes in  $Tree(r)$ 
    limited by  $M$ ;
```

Algorithm Description: Algorithm 1 consists of two subroutines. The first subroutine is to construct the user relationships as a tree (Lines 1-5). The tree is rooted by the single-relationship R_{single}^m and takes the group-relationship R_{group}^n as node, where the branches denote a user linking single-relationship and group-relationship (Line 5). The second subroutine is to anonymize the user PoIs in time-constrained data set \mathcal{D} (Lines 6-11). According to Section 3.2, the m prior knowledge is meaningless in group relationship if n is bigger than m .

Therefore, the user PoIs of single-relationship are only anonymized if $m \leq n$ (Lines 6-7), otherwise both user PoIs of single-relationship and group-relationship are anonymized (Lines 9- 11). The anonymous strategy is executed in a tree, iteratively (Line 11). If one node have been anonymized, its descendant nodes are verified whether there still exists group-relationship or not. If one descendant node still satisfies the group-relationship, the node is anonymized and verifies its descendant nodes.

Experimental Result and Conclusions

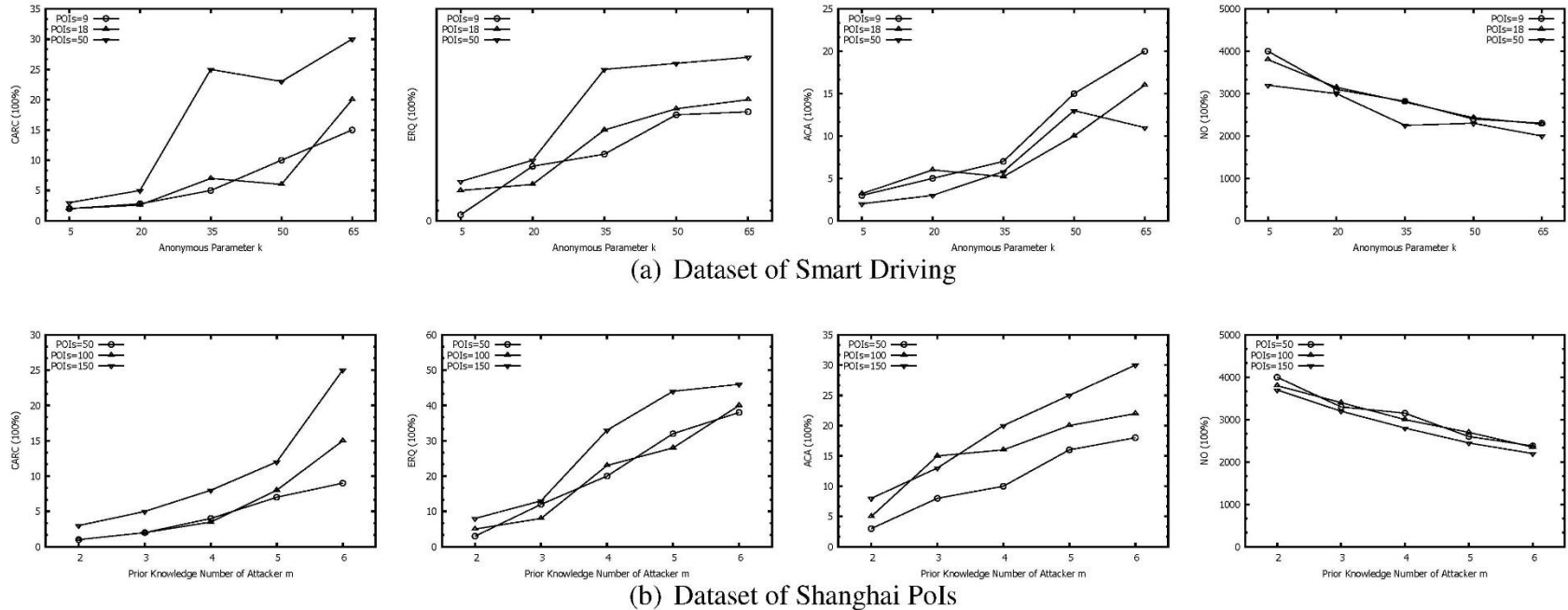


Fig.2: Experiments of CARC, ERQ, Average CA(ACA) and Number of changing PoIs

Experimental Result: Figure 2 shown that the values of CARC, ERQ and ACA appear an overall upward trend as the increase of m and the unmodified number of points of interest NO has been shown a downward trend. This upward trend mean that the anonymous data has a worse data and query availability and decreased number of NO refers to a more laborious anonymous operations. In general, the larger the value of m is, the worse the data availability of anonymous data is.

Conclusions: This paper designed an algorithm to protect the privacy relationship exposed by the user PoIs. Through the experimental evaluation, our method can protect the time-constrained PoIs under the hypothetical prior knowledge of attackers.